

General Instructions: Same as in Homework 1.

Honor Principle: Same as in Homework 1.

24. Recall our definition of the class $IP_{\alpha,\beta}$: We say that a language $L \subseteq \{0,1\}^*$ is in this class if there is a polynomial-time verifier V that uses a random string r and has the following properties, where P is an arbitrarily powerful prover that interacts with V :

$$\begin{aligned}x \notin L &\implies \forall P : \Pr_r[V * P(x, r) = 1] \leq \alpha, \\x \in L &\implies \exists P : \Pr_r[V * P(x, r) = 1] \geq \beta.\end{aligned}$$

We defined $IP = IP_{\frac{1}{3}, \frac{2}{3}}$ and remarked that the choice of the constants isn't terribly important, as can be proven by suitable repetition and Chernoff bound analysis. We also remarked that β can be made equal to 1 (perfect completeness), though not by simple repetition. Finally, we remarked that α cannot be made zero (perfect soundness), because that would boil the underlying class to plain old NP.

Justify this last remark. Specifically, prove that $IP_{0, \frac{2}{3}} = NP$. [2 points]

25. Let p be a prime. This problem involves the group \mathbb{Z}_p , consisting of integers $\{1, 2, \dots, p-1\}$ with multiplication performed mod p . At some point you will need to use the fact that every element of \mathbb{Z}_p has a multiplicative inverse mod p (that's what makes it a group).

The *quadratic residuosity problem* asks whether a given integer is a square mod p . The brute force solution is to try out all elements of \mathbb{Z}_p and compute the square of each, but it takes time proportional to p , which is exponential in the input length. But one can give interesting interactive proofs for this problem. To be precise, define the languages

$$\begin{aligned}QR &= \{ \langle p, x \rangle : p \text{ is prime, } x \in \mathbb{Z}_p, \text{ and } \exists y \in \mathbb{Z}_p (y^2 \equiv x \pmod{p}) \}, \\QNR &= \{ \langle p, x \rangle : p \text{ is prime, } x \in \mathbb{Z}_p, \text{ and } \forall y \in \mathbb{Z}_p (y^2 \not\equiv x \pmod{p}) \}.\end{aligned}$$

The acronyms denote "quadratic residue" and "quadratic non-residue," respectively.

Prove that both these languages are in IP and that one of these is in fact in NP. [2 points]

Hint: Your protocol for one of the languages should mimic the one we gave in class for GNI (graph non-isomorphism). Suppose $\langle p, x \rangle \in QNR$ and $z \in \mathbb{Z}_p$. What can you say about $xz^2 \pmod{p}$?